

# A "joint+marginal" approach to parametric optimization

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This work is based on:

## The "joint+marginal" approach to parametric polynomial optimization

To appear in **SIAM J. Optimization**

- Parametric Optimization
- The "joint+marginal" approach
- A hierarchy of semidefinite relaxations
- -> in optimization yields a "joint+marginal" algorithm

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# Parametric Optimization

Let  $\mathbf{Y} \subset \mathbb{R}^p$  be a compact set, called the **parameter** set.

Let  $\mathbf{K} \subset \mathbb{R}^n \times \mathbb{R}^p$  be the set:

$$\mathbf{K} := \{ (\mathbf{x}, \mathbf{y}) : \mathbf{y} \in \mathbf{Y}; \quad h_j(\mathbf{x}, \mathbf{y}) \geq 0, \quad j = 1, \dots, m \},$$

for some continuous functions  $h_j : \mathbb{R}^n \times \mathbb{R}^p \rightarrow \mathbb{R}$ .

Consider the following optimization problem:

$$J(\mathbf{y}) := \inf_{\mathbf{x}} \{ f(\mathbf{x}, \mathbf{y}) : \mathbf{x} \in \mathbf{K}_{\mathbf{y}} \},$$

where for each  $\mathbf{y} \in \mathbf{Y}$ , the  $\mathbf{K}_{\mathbf{y}} \subset \mathbb{R}^n$  is defined by:

$$\mathbf{K}_{\mathbf{y}} := \{ \mathbf{x} \in \mathbb{R}^n : (\mathbf{x}, \mathbf{y}) \in \mathbf{K} \}$$

## Parametric optimization is concerned with:

- the **global optimal value** function  $\mathbf{y} \mapsto J(\mathbf{y})$ , and
- the **global minimizer** set function  $\mathbf{y} \mapsto \mathbf{x}_j^*(\mathbf{y})$
- the **optimal dual multiplier** set function  $\mathbf{y} \mapsto \lambda_j^*(\mathbf{y})$  associated with the constraint  $h_j(\mathbf{x}, \mathbf{y}) \geq 0$ .

In general, getting **full** information is impossible, and one is satisfied with **local information** (e.g. **sensitivity analysis**) around some (even **local**) minimizer  $\mathbf{x}^*(\mathbf{y}) \in \mathbf{K}_{\mathbf{y}}$ ,  $\mathbf{y} \in \mathbf{Y}$ . (See e.g. the book by **Bonnans and Shapiro**.)

This talk

For **polynomial optimization** much more is possible!



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# The "joint+marginal" approach

Let  $\varphi$  be a Borel probability measure on  $\mathbf{Y}$ , with a positive density with respect to the Lebesgue measure on the smallest affine variety that contains  $\mathbf{Y}$ . For instance,

$$\varphi(B) := \left( \int_{\mathbf{Y}} d\mathbf{y} \right)^{-1} \int_B d\mathbf{y}, \quad \forall B \in \mathcal{B}(\mathbf{Y}),$$

is **uniformly distributed** on  $\mathbf{Y}$ .

For a **discrete** set of parameters  $\mathbf{Y}$  (finite or countable) take for  $\varphi$  a discrete probability measure on  $\mathbf{Y}$  with strictly positive weight at each point of the support.

Sometimes, e.g. in the context of **optimization with data uncertainty**,  $\varphi$  is already specified.

## A related infinite-dimensional linear program:

Consider the infinite-dimensional LP:

$$\mathbf{P} : \quad \rho := \inf_{\mu \in \mathbf{M}(\mathbf{K})} \left\{ \int_{\mathbf{K}} f d\mu : \pi\mu = \varphi \right\}$$

where:  $\mathbf{M}(\mathbf{K})$  is the set of Borel probability measures on  $\mathbf{K}$ , and  $\pi : \mathbf{M}(\mathbf{K}) \rightarrow \mathbf{M}(\mathbf{Y})$  is the **projection** (or, **marginal**) on  $\mathbf{Y}$ .

Whence the name "joint+marginal"-approach since:

- $\mu$  is a **joint distribution** on the variables  $\mathbf{x}$  **AND** the parameters  $\mathbf{y}$ .
- $\varphi$  is the **marginal** of  $\mu$  on  $\mathbf{Y}$  (fixed, as a constraint on  $\mu$ ).

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The dual  $\mathbf{P}^*$  is the infinite-dimensional LP:

$$\mathbf{P}^* : \quad \rho^* := \sup_{g \in C(\mathbf{Y})} \int_{\mathbf{Y}} g(\mathbf{y}) d\varphi(\mathbf{y})$$
$$f(\mathbf{x}, \mathbf{y}) - g(\mathbf{y}) \geq 0 \quad \forall (\mathbf{x}, \mathbf{y}) \in \mathbf{K}.$$

where  $C(\mathbf{Y})$  is the set of continuous functions on  $\mathbf{Y}$ .

In other words, among the continuous functions  $g$  on  $\mathbf{Y}$  such that:

$$f(\mathbf{x}, \mathbf{y}) \geq g(\mathbf{y}) \quad \forall \mathbf{x} \in \mathbf{K}_{\mathbf{y}},$$

one searches for the one that maximizes  $\int_{\mathbf{Y}} g d\varphi$ .

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# Why those LPs?

We assume that  $\mathbf{K}$  is compact.

As we shall see ....

Any optimal solution  $\mu^*$  of the primal  $\mathbf{P}$  encodes *all information* on the optimal solutions  $\mathbf{x}^*(\mathbf{y})$  of  $\mathbf{P}_{\mathbf{y}}$ .

Similarly ....

There is *no duality gap*  $\rho = \rho^*$  and so, in particular, the optimal value function  $\mathbf{y} \mapsto J(\mathbf{y})$  of  $\mathbf{P}_{\mathbf{y}}$  can be nicely *approximated by polynomials*.



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## Theorem (The primal side ...)

Assume that  $\mathbf{K}$  is compact and  $\mathbf{K}_{\mathbf{y}} \neq \emptyset$  for every  $\mathbf{y} \in \mathbf{Y}$ . Let

$\mathbf{X}_{\mathbf{y}}^* := \{\mathbf{x} \in \mathbb{R}^n : f(\mathbf{x}, \mathbf{y}) = J(\mathbf{y})\}$ ,  $\mathbf{y} \in \mathbf{Y}$ . Then:

(a)  $\rho = \int_{\mathbf{Y}} J(\mathbf{y}) d\varphi(\mathbf{y})$  and  $\mathbf{P}$  has an optimal solution.

(b) For every optimal solution  $\mu^*$  of  $\mathbf{P}$ , and for  $\varphi$ -almost all  $\mathbf{y} \in \mathbf{Y}$ , there is a probability measure  $\psi^*(d\mathbf{x} | \mathbf{y})$  on  $\mathbb{R}^n$ , concentrated on  $\mathbf{X}_{\mathbf{y}}^*$ , such that:

$$\mu^*(C \times B) = \int_B \psi^*(C | \mathbf{y}) d\varphi(\mathbf{y}), \quad \forall B \in \mathcal{B}(\mathbf{Y}), C \in \mathcal{B}(\mathbb{R}^n).$$

continued ...

(c) Assume that for  $\varphi$ -almost all  $\mathbf{y} \in \mathbf{Y}$ , the set of minimizers  $\mathbf{X}_{\mathbf{y}}^*$  is the singleton  $\{\mathbf{x}^*(\mathbf{y})\}$  for some  $\mathbf{x}^*(\mathbf{y}) \in \mathbf{K}_{\mathbf{y}}$ . Then there is a measurable mapping  $g : \mathbf{Y} \rightarrow \mathbf{K}_{\mathbf{y}}$  such that

$$g(\mathbf{y}) = \mathbf{x}^*(\mathbf{y}) \text{ for every } \mathbf{y} \in \mathbf{Y}; \quad \rho = \int_{\mathbf{Y}} f(g(\mathbf{y}), \mathbf{y}) d\varphi(\mathbf{y}),$$

and for every  $\alpha \in \mathbb{N}^n$ , and  $\beta \in \mathbb{N}^p$ :

$$\int_{\mathbf{K}} \mathbf{x}^{\alpha} \mathbf{y}^{\beta} d\mu^*(\mathbf{x}, \mathbf{y}) = \int_{\mathbf{Y}} \mathbf{y}^{\beta} g(\mathbf{y})^{\alpha} d\varphi(\mathbf{y}).$$

## Theorem (The dual side ...)

(a) There is *no duality gap*, i.e.,

$$\rho = \rho^* = \int_{\mathbf{Y}} J(\mathbf{y}) d\varphi(\mathbf{y}),$$

(b) One may use polynomials of  $\mathbb{R}[\mathbf{y}]$  to approximate  $\rho^*$ .

(c) Let  $(p_i) \subset \mathbb{R}[\mathbf{y}]$  be any *maximizing sequence*. Then:  
 *$L_1$ -norm convergence:*

$$\text{as } i \rightarrow \infty, \quad \int_{\mathbf{Y}} |J(\mathbf{y}) - p_i(\mathbf{y})| d\varphi(\mathbf{y}) \rightarrow 0$$

*$\varphi$ -almost sure convergence:* Let  $\tilde{p}_i := \max_{k=0,\dots,i} p_k$ . Then

$$\text{as } i \rightarrow \infty, \quad \tilde{p}_i \rightarrow J \quad \varphi\text{-almost surely in } \mathbf{Y}$$

# Polynomial Parametric Optimization

In general,  $P$  and  $P^*$  are **intractable!**

However .... when:

- $Y$  and  $K$ , are **basic semi-algebraic sets**, and:
- either one already knows the moments of  $\varphi$ , or  $Y$  is **simple enough** (e.g. a box, a simplex, a hyper-sphere) so that they can be computed.

... then one can approximate the optimal value  $\rho$  of  $P$ , and:

- The **optimal value mapping**  $\mathbf{y} \mapsto J(\mathbf{y})$
- The **global minimizer mapping**  $\mathbf{y} \mapsto \mathbf{x}^*(\mathbf{y})$ ,

... via the hierarchy of semidefinite relaxations

of the **moment-s.o.s. approach** in polynomial optimization.

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# The moment-s.o.s. approach

Let  $\mathbb{N}_i^n := \{\alpha \in \mathbb{N}^n : \sum_j \alpha_j \leq i\}$ .

With a sequence  $\mathbf{z} = (z_{\alpha\beta})$ , indexed in the canonical basis  $(\mathbf{x}^\alpha \mathbf{y}^\beta)$  of  $\mathbb{R}[\mathbf{x}, \mathbf{y}]$ , let  $L_{\mathbf{z}} : \mathbb{R}[\mathbf{x}, \mathbf{y}] \rightarrow \mathbb{R}$  be the linear mapping:

$$f (= \sum_{\alpha\beta} f_{\alpha\beta} \mathbf{x}^\alpha \mathbf{y}^\beta) \mapsto L_{\mathbf{z}}(f) := \sum_{\alpha\beta} f_{\alpha\beta} z_{\alpha\beta}, \quad f \in \mathbb{R}[\mathbf{x}, \mathbf{y}].$$

## The moment matrix $\mathbf{M}_i(\mathbf{z})$

associated with a sequence  $\mathbf{z} = (z_{\alpha\beta})$ , has its rows and columns indexed in the canonical basis  $(\mathbf{x}^\alpha \mathbf{y}^\beta)$ , and with entries.

$$\mathbf{M}_i(\mathbf{z})(\alpha, \beta), (\delta, \gamma)) = L_{\mathbf{z}}(\mathbf{x}^\alpha \mathbf{y}^\beta \mathbf{x}^\delta \mathbf{y}^\gamma) = z_{(\alpha+\delta)(\beta+\gamma)},$$

for every  $\alpha, \delta \in \mathbb{N}_i^n$  and every  $\beta, \gamma \in \mathbb{N}_i^p$



Let  $q$  be the polynomial  $(\mathbf{x}, \mathbf{y}) \mapsto q(\mathbf{x}, \mathbf{y}) := \sum_{u,v} q_{uv} \mathbf{x}^u \mathbf{y}^v$ .

### The localizing matrix $\mathbf{M}_i(q, \mathbf{z})$

associated with  $q \in \mathbb{R}[\mathbf{x}, \mathbf{y}]$  and a sequence  $\mathbf{z} = (z_{\alpha\beta})$ , has its rows and columns indexed in the canonical basis  $(\mathbf{x}^\alpha \mathbf{y}^\beta)$ , and with entries.

$$\begin{aligned} \mathbf{M}_i(q, \mathbf{z})(\alpha, \beta), (\delta, \gamma) &= L_{\mathbf{z}}(q(\mathbf{x}, \mathbf{y}) \mathbf{x}^\alpha \mathbf{y}^\beta \mathbf{x}^\delta \mathbf{y}^\gamma) \\ &= \sum_{u \in \mathbb{N}^n, v \in \mathbb{N}^p} q_{uv} z_{(\alpha+\delta+u)(\beta+\gamma+v)}, \end{aligned}$$

for every  $\alpha, \delta \in \mathbb{N}_i^n$  and every  $\beta, \gamma \in \mathbb{N}_i^p$ .

Let  $h_j \in \mathbb{R}[\mathbf{x}, \mathbf{y}]$  for every  $j = 1, \dots, m$ , and recall that

$$\mathbf{K} := \{(\mathbf{x}, \mathbf{y}) : \mathbf{y} \in \mathbf{Y}; \quad h_j(\mathbf{x}, \mathbf{y}) \geq 0, \quad j = 1, \dots, m\}.$$

The **parameter set**  $\mathbf{Y}$  is the basic semi-algebraic set

$$\mathbf{Y} := \{\mathbf{y} \in \mathbb{R}^p : h_k(\mathbf{y}) \geq 0, \quad k = m+1, \dots, t\}$$

Let  $\varphi$  be a probability measure on  $\mathbf{Y}$ ,

absolutely continuous with respect to the Lebesgue measure,  
and whose moments  $\gamma = (\gamma_\beta)$  with

$$\gamma_\beta = \int_{\mathbf{Y}} \mathbf{y}^\beta d\varphi(\mathbf{y}), \quad \forall \beta \in \mathbb{N}^p,$$

are all known.

# Primal semidefinite relaxations:

Let  $v_k := \lceil (\deg h_k)/2 \rceil$  for every  $k = 1, \dots, t$  and let  $i_0 := \max[\lceil (\deg f)/2 \rceil, \max_k v_k]$ .

For  $i \geq i_0$ , consider the semidefinite program:

$$\begin{aligned} \rho_i = & \inf_{\mathbf{z}} L_{\mathbf{z}}(f) \\ \text{s.t. } & \mathbf{M}_i(\mathbf{z}) \succeq 0 \\ & \mathbf{M}_{i-v_j}(h_j \mathbf{z}) \succeq 0, \quad j = 1, \dots, t \\ & L_{\mathbf{z}}(\mathbf{y}^\beta) = \gamma_\beta, \quad \forall \beta \in \mathbb{N}_i^p. \end{aligned}$$

which is a **relaxation** of  $\mathbf{P}$ , and

$$\rho_{i_0} \leq \dots \leq \rho_{i-1} \leq \rho_i \leq \dots \leq \rho.$$

# Dual semidefinite relaxation

The dual reads:

$$\rho_i^* = \sup_{g, (\sigma_i)} \int_{\mathbf{Y}} g(\mathbf{y}) d\varphi(\mathbf{y}) \quad (= \sum_{\beta} g_{\beta} \gamma_{\beta})$$

$$\text{s.t.} \quad f(\mathbf{x}, \mathbf{y}) - g(\mathbf{y}) = \sigma_0(\mathbf{x}, \mathbf{y}) + \sum_{j=1}^t \sigma_j(\mathbf{x}, \mathbf{y}) h_j(\mathbf{x}, \mathbf{y})$$

$$p \in \mathbb{R}[\mathbf{y}]; \sigma_j \in \Sigma[\mathbf{x}, \mathbf{y}], \quad j = 1, \dots, t$$

$$\deg p \leq 2i, \deg \sigma_j h_j \leq 2i, \quad j = 1, \dots, t$$

(Compare with  $\mathbf{P}^*$  where  $f(\mathbf{x}, \mathbf{y}) - g(\mathbf{y}) \geq 0$  for all  $\mathbf{x} \in \mathbf{K}_y$ ).

## Theorem (Results for the Primal **P**)

(a)  $\rho_i \uparrow \rho$  as  $i \rightarrow \infty$ .

(b) Let  $z^i$  be a nearly optimal solution, e.g. such that  $L_{z^i}(f) \leq \rho_i + 1/i$ . If for  $\varphi$ -almost all  $y \in Y$ ,  $J(y)$  is attained at a unique optimal solution  $x^*(y)$ , then:

$$\lim_{i \rightarrow \infty} z^i_{\alpha\beta} = \int_Y y^\beta x^*(y)^\alpha d\varphi(y), \quad \forall \alpha \in \mathbb{N}^n, \beta \in \mathbb{N}^p.$$

In particular, for every  $k = 1, \dots, n$ ,

$$\lim_{i \rightarrow \infty} z^i_{e(k)\beta} = \int_Y y^\beta x_k^*(y) d\varphi(y), \quad \forall \beta \in \mathbb{N}^p,$$

where  $e(k)_j = \delta_{j=k}$ ,  $j = 1, \dots, n$  (with  $\delta$  being the Kronecker symbol).

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Assume  $\mathbf{x}_k^*(\mathbf{y}) \geq 0$  on  $\mathbf{Y}$ .

Then for sufficiently large  $i$ ,

$$\mathbf{z}^i_{e(k)\beta} \approx \int_{\mathbf{Y}} \mathbf{y}^\beta \mathbf{x}_k^*(\mathbf{y}) d\varphi(\mathbf{y}), \quad \forall \beta \in \mathbb{N}^p.$$

That is, one has a good approximation of all **moments** of the measure  $\mathbf{x}_k^*(\mathbf{y}) d\varphi(\mathbf{y})$  with **density**  $\mathbf{x}_k^*(\mathbf{y})$  on  $\mathbf{Y}$ .

Hence one may **approximate** the optimal  $k$ -th coordinate function  $\mathbf{y} \mapsto \mathbf{x}_k^*(\mathbf{y})$  by e.g. **maximum-entropy** estimation methods.



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## Theorem (Results for the dual $\mathbf{P}^*$ )

Consider the *dual semidefinite relaxations*. Then:

(a)  $\rho_i^* \uparrow \rho$  as  $i \rightarrow \infty$ .

(b) Let  $(\rho_i, (\sigma_j^i))$  be a nearly optimal solution e.g.. such that

$\int_{\mathbf{Y}} \rho_i d\varphi \geq \rho_i^* - 1/i$ . Then  $\rho_i \leq J(\cdot)$  and

$$\lim_{i \rightarrow \infty} \int_{\mathbf{Y}} |J(\mathbf{y}) - \rho_i(\mathbf{y})| d\varphi(\mathbf{y}) = 0$$

Moreover if one defines

$$\tilde{\rho}_0 := \rho_0, \quad \mathbf{y} \mapsto \tilde{\rho}_i(\mathbf{y}) := \max[\tilde{\rho}_{i-1}(\mathbf{y}), \rho_i(\mathbf{y})], \quad i = 1, 2, \dots,$$

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$\int_{\mathbf{Y}} \rho_i d\varphi \geq \rho_i^* - 1/i$ . Then  $\rho_i \leq J(\cdot)$  and

$$\lim_{i \rightarrow \infty} \int_{\mathbf{Y}} |J(\mathbf{y}) - \rho_i(\mathbf{y})| d\varphi(\mathbf{y}) = 0$$

Moreover if one defines

$$\tilde{\rho}_0 := \rho_0, \quad \mathbf{y} \mapsto \tilde{\rho}_i(\mathbf{y}) := \max[\tilde{\rho}_{i-1}(\mathbf{y}), \rho_i(\mathbf{y})], \quad i = 1, 2, \dots,$$

then  $\tilde{\rho}_i \rightarrow J(\cdot)$   $\varphi$ -almost uniformly on  $\mathbf{Y}$ .

# Illustrative Examples

With  $\mathbf{Y} := [0, 1]$ , let  $\mathbf{K} := \{(\mathbf{x}, \mathbf{y}) : 1 - x_1^2 - x_2^2 \geq 0\} \subset \mathbb{R}^2$ , and  $f(\mathbf{x}, \mathbf{y}) := \mathbf{y}x_1 + (1 - \mathbf{y})x_2$ . One easily obtains:

$$J(\mathbf{y}) = -\sqrt{\mathbf{y}^2 + (1 - \mathbf{y})^2},$$

and

$$\mathbf{x}_1^*(\mathbf{y}) = \frac{-\mathbf{y}}{\sqrt{\mathbf{y}^2 + (1 - \mathbf{y})^2}}; \quad \mathbf{x}_2^*(\mathbf{y}) = \frac{\mathbf{y} - 1}{\sqrt{\mathbf{y}^2 + (1 - \mathbf{y})^2}}.$$

With 8 moments of the uniform distribution of  $[0, 1]$ , one obtains:

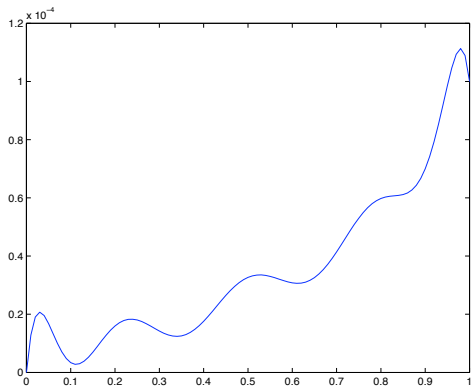


Figure:  $J(\mathbf{y}) - p_4(\mathbf{y})$  on  $[0, 1]$ . (Scale is  $10^{-4}$ )

And the Boltzmann-Shannon **maximum-entropy** estimation of  $\mathbf{x}^*(\mathbf{y})$  with **8 moments** gives:

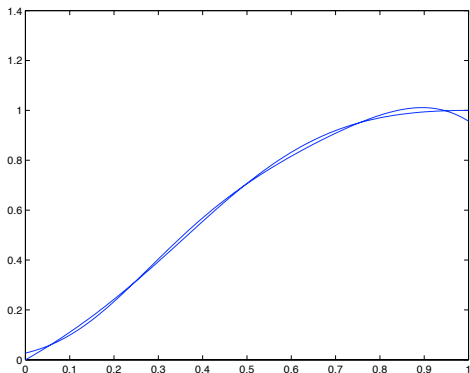


Figure: max-entropy estimate  $g_1(\mathbf{y})$  vs  $-\mathbf{x}^*(\mathbf{y}) = \mathbf{y}/\sqrt{\mathbf{y}^2 + (1 - \mathbf{y})^2}$

## Example 2:

$\mathbf{Y} = [0, 1]$ ,  $f(\mathbf{x}, \mathbf{y}) := (1 - 2\mathbf{y})(x_1 + x_2)$ , and

$$\mathbf{K} := \{(\mathbf{x}, \mathbf{y}) : \mathbf{y}x_1^2 + x_2^2 - \mathbf{y} \leq 0; x_1^2 + \mathbf{y}x_2^2 - \mathbf{y} \leq 0\}.$$

That is, for each  $\mathbf{y} \in \mathbf{Y}$  the set  $\mathbf{K}_{\mathbf{y}}$  is the intersection of two ellipsoids.  $J(\mathbf{y}) = -2|1 - 2\mathbf{y}| \sqrt{\frac{\mathbf{y}}{1+\mathbf{y}}}$ .

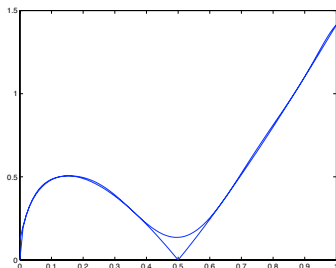


Figure:  $-p_2(\mathbf{y})$  and  $-J(\mathbf{y})$  on  $[0, 1]$  (4 moments)



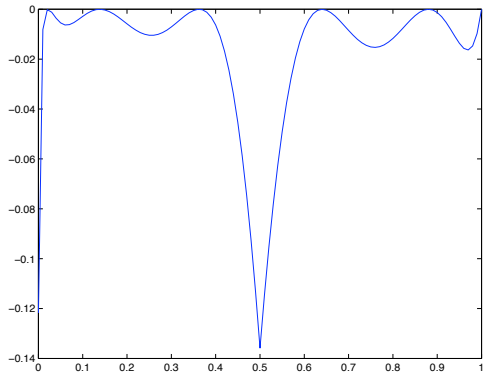


Figure: The curve  $p_2(\mathbf{y}) - J(\mathbf{y})$  on  $[0, 1]$  (4 moments)

## Example 3:

Consider the following system of 4 quadratic equations in 4 variables and one parameter  $y \in \mathbf{Y} = [0, 1]$ :

$$\begin{aligned}x_1 x_2 - x_1 x_3 - x_4 &= \mathbf{y} & x_2 x_3 - x_2 x_4 - x_1 &= \mathbf{y} \\ -x_1 x_3 + x_3 x_4 - x_2 &= \mathbf{y} & x_1 x_4 - x_2 x_4 - x_3 &= \mathbf{y},\end{aligned}$$

for which one wishes to compute the **minimum norm**  $J(\mathbf{y})$  of **real solutions** as a function of  $\mathbf{y} \in \mathbf{Y}$ .

	$\mathbf{y}=0.1$	$\mathbf{y}=0.5$	$\mathbf{y}=1$
$J(\mathbf{y})$	0.0400	1.0000	2.0000
$p_6(\mathbf{y})$	0.0384	0.9264	1.9887
$p_8(\mathbf{y})$	0.0390	0.9395	2.0000

Table:  $J(\mathbf{y})$  versus  $p_k(\mathbf{y})$

# A "joint+marginal" algorithm for Optimization

Given  $\mathbf{K} \subset \mathbb{R}^n$  and  $f \in \mathbb{R}[\mathbf{x}]$ , consider the polynomial optimization problem  $\mathbf{P} : \min_{\mathbf{x}} \{f(\mathbf{x}) : \mathbf{x} \in \mathbf{K}\}$ .

IDEA: Consider  $\mathbf{x}_1$  as a **parameter**  $\mathbf{y}$  in some interval  $\mathbf{Y} \subset \mathbb{R}$  (to be determined, e.g., easily when  $\mathbf{K}$  is convex), so that:

$$J(\mathbf{y}) = \min_{\mathbf{x}} \{f(\mathbf{x}) : \mathbf{x} \in \mathbf{K}; \mathbf{x}_1 = \mathbf{y}\}, \quad \mathbf{y} \in \mathbf{Y}.$$

- Compute a **polynomial approximation**  $p_k(\mathbf{y})$  of  $J(\mathbf{y})$  via the  $k$ -th **semidefinite relaxation** of the "joint+marginal" approach.
- Minimize the **univariate** polynomial  $p_k(\mathbf{y})$  on  $\mathbf{Y}$  (easy! reduces to solving a **single** semidefinite program), and get  $z_1 \in \mathbf{Y}$ .
- In  $\mathbf{P}$  fix  $\mathbf{x}_1 := z_1$ , and repeat for a  $(n-1)$  optimization problem to obtain  $\mathbf{x}_2 = z_2$ , etc. until we get  $(z_1, z_2, \dots, z_n)$

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## The rationale behind the "joint+marginal" algorithm:

- The larger  $k$ , the better the approximation of  $J(\mathbf{y})$  by the univariate polynomial  $p_k(\mathbf{y})$ . And so in minimizing  $p_k(\mathbf{y})$  over  $\mathbf{Y}$  one has a good chance to obtain  $z_1 \approx \mathbf{x}_1^*$ , where  $\mathbf{x}^*$  is a global minimizer of  $\mathbf{P}$ . And so at the end one may expect  $z \approx \mathbf{x}^*$ .
- The interest is to precisely have  $k$  not too large so as to handle relatively large size problems, and use the output point  $z$  as the initial point of a local minimization algorithm to next obtain a local minimizer  $\tilde{\mathbf{x}} \in \mathbf{K}$  with reasonable hope that  $\tilde{\mathbf{x}}$  is not far away of  $\mathbf{x}^*$ .
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**THANK YOU !!**